

DEF. VON NEUMANN ENTROPY

$$S(p) := -\text{tr}(p \log p) \quad \text{where } p \text{ is a density matrix}$$

$$P = \sum_{i=1}^d p_i |v_i\rangle \langle v_i| \implies S(p) = -\sum_{i=1}^d p_i \log(p_i) = H(\sum p_i)$$

↑  
EIGENDEC.  
eigenvalues  
(with  $\sum_{i=1}^d p_i = 1$ )

↑  
SHANNON  
ENTROPY

PROOF:

$$\text{In general we have } f(p) = \sum_{i=1}^d f(p_i) |v_i\rangle \langle v_i|.$$

FACT  $\oplus$ :

This because every  $f(\cdot)$  can be written as  $f(p) = \sum_{s=0}^{\infty} c_s p^s$

$$\begin{aligned} \Rightarrow f(p) &= \sum_{s=0}^{\infty} c_s p^s \\ &= \sum_{s=0}^{\infty} c_s \underbrace{\sum_{i=1}^d |v_i\rangle \langle v_i|}_{\sum_{i=1}^d |v_i\rangle \langle v_i|} = \sum_{s=0}^{\infty} \sum_{i=1}^d c_s \underbrace{(p^s |v_i\rangle)}_{p_i^s |v_i\rangle} \langle v_i| = \\ &= \sum_i \left( \sum_s c_s p_i^s \right) |v_i\rangle \langle v_i| = \sum_{i=1}^d \underbrace{f(p_i)}_{\sum_s c_s p_i^s} |v_i\rangle \langle v_i| \end{aligned}$$

$$\therefore \text{given } f(p) := p \log(p) \Rightarrow p \log(p) = \sum_{i=1}^d (p_i \log(p_i)) |v_i\rangle \langle v_i|$$

$$\begin{aligned} \Rightarrow S(p) &:= -\text{tr}(p \log p) = -\sum_{i=1}^d (p_i \log(p_i)) \underbrace{\text{tr}(|v_i\rangle \langle v_i|)}_{1} \\ &= -\sum_{i=1}^d p_i \log p_i \end{aligned}$$

QBS.

$$0 \leq S(p) \leq \log(d)$$

PROOF:

$$S(p) = H(\sum p_i) \quad \text{and we showed that } 0 \leq H(x) \leq \log(d).$$

OBS)

$$\cdot S(p) = 0 \iff p = |\psi\rangle\langle\psi| \quad (\text{PURE})$$

$$\cdot S(p) = \log(d) \iff p = \frac{1}{d} \quad (\text{MAXIMALLY MIXED STATE})$$

PROOF:

$$\cdot S(p) = H(\sum_i p_i \hat{P}_i) = 0 \iff \{\hat{P}_i\}_{i=1}^d \text{ is a peaked distribution i.e. } p_i = \delta_{i,a} \quad \forall i.$$

↑  
PREVIOUS  
LECTURE

$$\cdot S(p) = H(\sum_i p_i \hat{P}_i) = \log(d) \iff \{\hat{P}_i\}_{i=1}^d \text{ is a flat distribution i.e. } p_i = \frac{1}{d} \quad \forall i.$$

$$\Leftrightarrow p = \sum_{i=1}^d p_i |\psi_i\rangle\langle\psi_i| = \frac{1}{d} \underbrace{\sum_{i=1}^d |\psi_i\rangle\langle\psi_i|}_{\mathbb{I}} \Rightarrow \mathbb{I}$$

$$\cdot S(p) = S(U_p U^\dagger) \quad \forall U \text{ unitary.}$$

PROOF:

$P = \sum_i p_i |\psi_i\rangle\langle\psi_i|$  and  $(U_p U^\dagger) = \sum_i p_i (U|\psi_i\rangle\langle\psi_i|U^\dagger)$  have the same eigenvalues.

and  $S(\cdot)$  depends only by the eigenvalues.

• DEF: (RELATIVE ENTROPY)

$$S(P_1 \| P_2) := \text{tr} [P_1 \log(P_2)] - \text{tr} [P_1 \log(P_1)]$$

$$= -S(P_1) - \text{tr} [P_1 \log(P_2)]$$

THEOREM: (KLEIN INEQUALITY)

$$S(P_1 \parallel P_2) \geq 0 \quad \left( S(P_1 \parallel P_2) = 0 \iff P_1 = P_2 \right)$$

PROOF:

$$P_1 = \sum_{i=1}^d p_i |V_i\rangle\langle V_i|, \quad P_2 = \sum_{j=1}^d q_j |W_j\rangle\langle W_j|$$

$$\begin{aligned} S(P_1 \parallel P_2) &= -S(P_1) - \text{tr}[P_1 \log(P_2)] \\ &= -S(P_1) - \text{tr}\left[\sum_{i=1}^d p_i |V_i\rangle\langle V_i| \log(P_2)\right] \end{aligned}$$

$$= -S(P_1) - \sum_{i=1}^d p_i \langle V_i | \log(P_2) | V_i \rangle$$

$$\begin{aligned} &= -S(P_1) - \sum_{i=1}^d \sum_{j=1}^d p_i \log(q_j) \underbrace{\langle V_i | W_j \rangle \langle W_j | V_i \rangle}_{\|V_i\| \|W_j\|^2} = \\ \log(P_2) &= \sum_{j=1}^d \log(q_j) |W_j\rangle\langle W_j| \quad \text{with } |W_j\rangle^2 \end{aligned}$$

★

$$= -S(P_1) - \sum_{i=1}^d \sum_{j=1}^d p_i \log(q_j) |V_i\langle W_j\rangle|^2$$

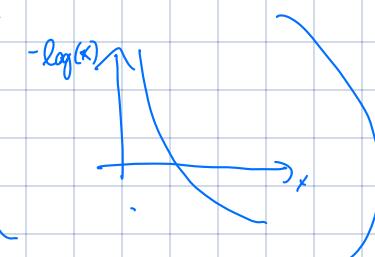
$$= + \sum_{i=1}^d p_i \log(p_i) - \sum_{i=1}^d \sum_{j=1}^d p_i \log(q_j) |V_i\langle W_j\rangle|^2$$

$$= + \sum_{i=1}^d \sum_{j=1}^d |V_i\langle W_j\rangle|^2 p_i \log(p_i) - \sum_{i=1}^d \sum_{j=1}^d |V_i\langle W_j\rangle|^2 p_i \log(q_j)$$

$$\sum_j |V_i\langle W_j\rangle|^2 = 1$$

$$= + \sum_{i=1}^d \sum_{j=1}^d |V_i\langle W_j\rangle|^2 p_i \left( \log\left(\frac{p_i}{q_j}\right) \right) = \sum_{i=1}^d \sum_{j=1}^d |V_i\langle W_j\rangle|^2 p_i \left( -\log\left(\frac{q_j}{p_i}\right) \right)$$

$$\geq \sum_i p_i \left[ -\log \left( \sum_j P(i,j) \frac{q_j}{p_i} \right) \right] = \sum_i p_i \left[ -\log \left( \frac{k_i}{p_i} \right) \right]$$

•  $\sum_j P(i,j) = 1$ ,  $P(i,j) \geq 0$   
 •  $-\log(x)$  CONVEX  


•  $k_i := \sum_j P(i,j) q_j$   
 { Prob. dist.  
 ( $\sum_i k_i = 1$ )  
 $k_i \geq 0$ )

$$= \sum_i p_i \log \left( \frac{p_i}{k_i} \right) \geq 0$$

↑  
Bochner's lecture

- The first inequality is saturated when  $p_i = q_j \forall i, j$   
 $\Downarrow$   
 $p_i = q_j = \frac{1}{d} \forall i, j$

or when  $P(i,j) = \delta_{j,i_0}$ .  $\forall j = 1, \dots, d$

- But  $P(i,j) = |\langle v_i | w_j \rangle|^2 = \delta_{j,i_0} \forall j = 1, \dots, d \Leftrightarrow |v_i\rangle = |w_{i_0}\rangle \forall i$

$$|\langle v_i \rangle| = \sum_j |\langle w_j | v_i \rangle| |\langle w_j \rangle|$$

$\uparrow$   
 $\delta_{j,i_0}$  &  $\langle w_j \rangle$   
 $= \langle w_{i_0} \rangle$

- The second inequality is saturated when  $p_i = k_i \forall i$  (previous lecture)

$$\Leftrightarrow p_i = \sum_j P(i,j) q_j = \sum_j \delta_{j,i_0} q_j = q_{i_0} \forall i$$

$$\Leftrightarrow p_1 = \sum_i p_i |\langle v_i | \rangle \langle v_i | = \sum_i q_{i_0} |\langle w_{i_0} | \rangle \langle w_{i_0} | = p_2$$

- OBS.
    - $S(P_1 \parallel P_2) \neq S(P_2 \parallel P_1)$  (NOT SYMMETRIC)
    - $\text{Supp}(P_1) \cap \text{Ker}(P_2) \neq \emptyset \Rightarrow S(P_1 \parallel P_2) = +\infty$
- PROOF:
- $$S(P_1 \parallel P_2) = -S(P_1) - \sum_{i=1}^d \sum_{j=1}^d p_i \log(p_j) |\langle v_i | w_j \rangle|^2$$
- If  $p_j = 0$  but  $p_i |\langle v_i | w_j \rangle| \neq 0 \quad \forall i$   $\Rightarrow S(P_1 \parallel P_2) = +\infty$
- $\Updownarrow$
- $w_j \notin \text{Ker}(P_2) \quad \text{if } w_j \in \text{Ker}(P_2) = \text{Supp}(P_2)$
- So  $\text{Ker}(P_2) \cap \text{Supp}(P_1) \neq \emptyset \Rightarrow S(P_1 \parallel P_2) = +\infty \Rightarrow \square$
- 

•  $O = O_1 \otimes O_2$  with  $O_1 = O_1^+$ ,  $O_2 = O_2^+$

$$\log(O) = \log(O_1) \otimes \mathbb{1} + \mathbb{1} \otimes \log(O_2) = \star$$

PROOF:

$$\log(O) = \log(O_1 \otimes O_2) = \log \left[ \sum_{i,s} \lambda_i^{(1)} \lambda_s^{(2)} \left( |v_i^{(1)}\rangle \langle v_i^{(1)}| \otimes |v_s^{(2)}\rangle \langle v_s^{(2)}| \right) \right] =$$

$O_1 = \sum_i \lambda_i^{(1)} |v_i^{(1)}\rangle \langle v_i^{(1)}|$  eigenvectors for  $O_1 \otimes O_2$

$O_2 = \sum_s \lambda_s^{(2)} |v_s^{(2)}\rangle \langle v_s^{(2)}|$

$$= \sum_{i,s} \log(\lambda_i^{(1)} \lambda_s^{(2)}) \left( |v_i^{(1)}\rangle \langle v_i^{(1)}| \otimes |v_s^{(2)}\rangle \langle v_s^{(2)}| \right) =$$

$f(A) = \sum_i f(\lambda_i) |v_i\rangle \langle v_i|$

$\sum_i |v_i\rangle \langle v_i|$

$\log(ab) = \log(a) + \log(b)$

$$= \sum_{i,s} \log(\lambda_i^{(1)}) \left( |v_i^{(1)}\rangle \langle v_i^{(1)}| \otimes |v_s^{(2)}\rangle \langle v_s^{(2)}| \right) + \sum_{i,s} \log(\lambda_s^{(2)}) \left( |v_i^{(1)}\rangle \langle v_i^{(1)}| \otimes |v_s^{(2)}\rangle \langle v_s^{(2)}| \right)$$

$$= \sum_i \log(\lambda_i^{(1)}) \left( |v_i^{(1)}\rangle \langle v_i^{(1)}| \otimes \mathbb{1} \right) + \sum_s \log(\lambda_s^{(2)}) \left( \mathbb{1} \otimes |v_s^{(2)}\rangle \langle v_s^{(2)}| \right) =$$

$$= \log(Q_1) \otimes \mathbb{1} + \log(Q_2) \otimes \mathbb{1}$$

COR.

$$\cdot Q = Q_1 \otimes Q_2 \otimes \cdots \otimes Q_n, \quad Q_i = Q_i^+ \quad \forall i$$

$$\cdot \log(Q) = \sum_{i=1}^N \mathbb{1}_{i,i} \otimes \log(Q_i) \otimes \mathbb{1}_{i+1,\dots,n}$$

COR:

$$S(P_1 \otimes P_2 \otimes \cdots \otimes P_N) = \sum_{i=1}^N S(P_i)$$

PROOF:

$$\log(P) = \sum_{i=1}^N \mathbb{1}_{i,i} \otimes \log(p_i) \otimes \mathbb{1}_{i+1,\dots,n}$$

$$\begin{aligned} S(P) &= -\text{tr}_2\left(P \log(P)\right) = -\sum_{i=1}^N \text{tr}_2\left(P \left(\mathbb{1}_{i,i} \otimes \log(p_i) \otimes \mathbb{1}_{i+1,\dots,n}\right)\right) = \\ &= -\sum_{i=1}^N \text{tr}_2\left(\underset{P_i}{\text{tr}_i}(P) \log(p_i)\right) = \sum_{i=1}^N S(P_i) \end{aligned}$$

THEOR. ②

$$\cdot P_A = \text{tr}_B(P_{AB}), \quad P_B = \text{tr}_A(P_{AB})$$

$$\cdot S(P_{AB}) \leq S(P_A) + S(P_B) \quad \left(S(P_{AB}) = S(P_A) + S(P_B) \iff P_{AB} = P_A \otimes P_B\right)$$

PROOF:

$$\cdot S(P_{AB} \parallel P_A \otimes P_B) \geq 0$$

↑  
KLEIN INEQUALITY

$$\cdot \underbrace{S(P_{AB} \parallel P_A \otimes P_B)}_{\geq 0} = -S(P_{AB}) - \text{tr}_2\left(P_{AB} \frac{\log(P_A \otimes P_B)}{\mathbb{1} \otimes \mathbb{1}}\right)$$

$$\log(P_A) \otimes \mathbb{1} + \mathbb{1} \otimes \log(P_B)$$

$$= -S(P_{AB}) - \underbrace{\text{tr}(P_{AB} \log(P_A \otimes I))}_{\text{tr}(P_A \log(P_A))} - \text{tr}(P_{AB} \# \otimes \log(I_B))$$

$$= -S(P_{AB}) + S(P_A) + S(P_B) \geq 0 \Rightarrow \underline{\Sigma}$$

$$\stackrel{*}{=} \Leftrightarrow S(P_{AB} \parallel P_A \otimes P_B) = 0 \Leftrightarrow P_{AB} = P_A \otimes P_B$$

FACT ]  
③ :

$$P_A := \text{tr}_B(p), \quad P_B := \text{tr}_A(p)$$

$$\bullet \quad P = |\psi\rangle\langle\psi| \quad \text{pure} \Rightarrow S(P_A) = S(P_B)$$

(This is not tree in general if  $P$  is not prime)  
 Find a counter-example (5).

## PROOF:

SCHMIDT.

$$|n\rangle = \sum_{i=1}^{\min(d_A, d_B)} |\lambda_i| |V^A_i\rangle \otimes |V^B_i\rangle$$

### SUBPROOF:

$$\text{Nitch } \langle \sqrt{i}, \sqrt{s} \rangle = \delta_{i,s} , \quad \langle \sqrt[3]{i}, \sqrt[3]{s} \rangle = \delta_{i,s} \dots$$

$$\cdot W = \sum_{i,s} \langle i, s | W | i, s \rangle = \sum_{i,s} (U D V^+)^{i,s} | i, s \rangle =$$

$\uparrow$   
 $(C)_{i,s}$

SVD

$C = U D V^+$

$$= \sum_{\ell} \sum_{i \in S} U_{i,\ell} D_{\ell\ell}(V^+)_{\ell,i} |i\rangle \otimes |\ell\rangle = \sum_{\ell} \lambda_{\ell} \left( \sum_i U_{i,\ell} |i\rangle \right) \otimes \underbrace{\left[ \sum_{\ell} \left( V^+ \right)_{\ell,i} |\ell\rangle \right]}_{\begin{array}{c} \text{if } |\ell\rangle \\ \text{if } |i\rangle \end{array}}$$

$$Tr_B(|\psi\rangle\langle\psi|) = \sum_i \lambda_i |\psi_i^{(A)}\rangle\langle\psi_i^{(A)}|$$

$\Rightarrow$  Some eigenvalues  $\Rightarrow$  same entropies.

$$\text{tr}_A(|\psi\rangle\langle\psi|) = \sum_i \lambda_i |\psi_i^{(B)}\rangle\langle\psi_i^{(B)}|$$

## THEOR | ④

$$S(P_{AB}) \geq |S(P_A) - S(P_B)|$$

PROOF:

- Given a mixed state  $P_{AB}$ , then  $\exists R$  such that  $P_{AB} = \text{tr}_R [ |\psi\rangle\langle\psi|]$

with  $|\psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_R$

SUBPROOF (ALREADY SEEN)

E.g. take  $|\psi\rangle := \sqrt{P_{AB}} \otimes \frac{1}{\sqrt{d}} \left( \sum_{i=1}^d |i\rangle_{AB} \otimes |i\rangle_R \right)$  and verify.

$$\left( -\boxed{N_{AB}} \circ \boxed{\sqrt{P_{AB}}} = -\boxed{P_{AB}} \right)$$

- We have:

$$S(P_B) = S(P_{AR}) \stackrel{\substack{\uparrow \\ (\text{TH ②})}}{\leq} S(P_A) + S(P_R) = S(P_A) + S(P_{AB})$$

$$\begin{cases} \cdot S(P_B) = S(P_{AR}) \\ \quad \uparrow \\ \cdot \text{FACT ③} \\ \cdot P_{ABR} \text{ pure} \end{cases}$$

$$\begin{cases} \cdot S(P_{AB}) = S(P_R) \\ \quad \uparrow \\ \cdot \text{FACT ③} \\ \cdot P_{ABR} \text{ pure} \end{cases}$$

$$\Rightarrow \boxed{S(P_{AB}) \geq S(P_B) - S(P_A)} = \star_1$$

$$S(P_A) = S(P_{BR}) \stackrel{\substack{\uparrow \\ \text{FACT ③}}}{\leq} S(P_B) + S(P_R) = S(P_B) + S(P_{AB})$$

$$\Rightarrow \boxed{S(P_{AB}) \geq S(P_A) - S(P_B)} = \star_2$$

$\Rightarrow \square$   
 $\star_1, \star_2$

OBS] Also in the classical case I have  $H(X, Y) \geq |H(X) - H(Y)|$

But we did not mention it because in the classical case

We have an even stronger inequality:  $H(X, Y) \geq \max(H(X), H(Y))$

Which can be violated quantumly.

$$S(P_{AB}) \geq \max(S(P_A), S(P_B))$$

No!

" " 2      " 1      No!      No!

" " 1      " 0      FALSE!

PROOF.

COUNTER-EXAMPLE:  $P_{AB} = |\psi\rangle\langle\psi|$  with  $|\psi\rangle = |\downarrow\rangle$  max. entangled state.

$$\bullet S(P_{AB}) = 0$$

↑  
state pure.

$$\bullet \max(S(P_A), S(P_B)) = \max\left(S\left(\frac{1}{d_A}\right), S\left(\frac{1}{d_B}\right)\right) \neq 0$$

$\frac{1}{d_A}$        $\frac{1}{d_B}$

DEF.

$$P \text{ Separable} \iff P = \sum_{i=1}^N p_i P_1^{(i)} \otimes P_2^{(i)}$$

with  $\begin{cases} p_i > 0 \\ \sum_{i=1}^N p_i = 1 \end{cases}$

$\cdot P_1^{(i)}$  and  $P_2^{(i)}$  states.

DEF

$P$  entangled  $\iff P$  not separable.

FACT:

$$P \text{ Separable} \Rightarrow S(P_{AB}) \geq S(P_A)$$

PROOF: NOT SHOWN HERE.

# THEOR. (CONCAVITY)

$$S\left(\sum_{i=1}^N p_i P_i\right) \geq \sum_{i=1}^N p_i S(P_i)$$

$$\left( S\left(\sum_{i=1}^N p_i P_i\right) = \sum_{i=1}^N p_i S(P_i) \Leftrightarrow p_i = p \forall i \right)$$

PROOF:

$$P_{AB} := \sum_{i=1}^N p_i P_i \otimes |i\rangle\langle i| = \sum_{i=1}^N \sum_{s=1}^d p_i \lambda_s^{(i)} |V_s^{(i)}\rangle\langle V_s^{(i)}| \otimes |i\rangle\langle i|$$

↑  
EIGEN.

$$S(P_{AB}) = H\left(\{p_i \lambda_s^{(i)}\}\right) = -\sum_{i,s} p_i \lambda_s^{(i)} \log(p_i \lambda_s^{(i)}) = -\sum_{i,s} p_i \lambda_s^{(i)} \log(p_i) - \sum_{i,s} p_i \lambda_s^{(i)} \log(\lambda_s^{(i)})$$

$$= H\left(\{\sum_i p_i\}\right) + \sum_{i=1}^N p_i S(p_i)$$

↑  
 $(\sum_s \lambda_s^{(i)} = 1)$

$$P_A = \text{tr}_B(P_{AB}) = \sum_{i=1}^N p_i P_i$$

$$S(P_A) = S\left(\sum_{i=1}^N p_i P_i\right)$$

$$P_B = \text{tr}_A(P_{AB}) = \sum_{i=1}^N p_i |i\rangle\langle i|$$

$$S(P_B) = H\left(\{p_i\}\right)$$

$$S(P_{AB}) \leq S(P_A) + S(P_B)$$

↑  
SUB-ADD.

$$\underbrace{H\left(\{\sum_i p_i\}\right) + \sum_{i=1}^N p_i S(p_i)}_{\text{Left side}} \leq \underbrace{S\left(\sum_{i=1}^N p_i P_i\right) + H\left(\{\sum_i p_i\}\right)}_{\text{Right side}} \Rightarrow \square$$

THEOR.

$$\sum_{i=1}^N p_i S(p_i) \leq S\left(\sum_{i=1}^N p_i P_i\right) \leq H\left(\{\sum_i p_i\}\right) + \sum_{i=1}^N p_i S(p_i)$$

PROOF:

$$S\left(\sum_{i=1}^N p_i \rho_i\right) = S\left(\sum_{i=1}^N \sum_{j=1}^d p_i \lambda_j^{(i)} |v_j^{(i)}\rangle\langle v_j^{(i)}|\right)$$

$p_i = \sum_{j=1}^d \lambda_j^{(i)} |v_j^{(i)}\rangle\langle v_j^{(i)}|$

orthogonal basis

$$|\psi\rangle_{AB} := \sum_{i=1}^N \sum_{j=1}^d \sqrt{p_i \lambda_j^{(i)}} |v_j^{(i)}\rangle \otimes |(i,j)\rangle$$

$$S(P_A) = S(\text{tr}_B [|\psi\rangle\langle\psi|]) = S(\text{tr}_A [|\psi\rangle\langle\psi|]) = S(P_B)$$

||  $|\psi\rangle\langle\psi|$  pure

$$S\left(\sum_{i,j} p_i \lambda_j^{(i)} |v_j^{(i)}\rangle\langle v_j^{(i)}|\right)$$

||

$$S\left(\sum_{i=1}^N p_i \rho_i\right)$$

$$\Rightarrow S(P_B) = S\left(\sum_{i=1}^N p_i \rho_i\right)$$

$$\begin{aligned} P_B &= \text{tr}_A [|\psi\rangle\langle\psi|] = \text{tr}_A \left( \sum_{i=1}^N \sum_{j=1}^d \sum_{k=1}^N \sum_{l=1}^d \sqrt{p_i \lambda_j^{(i)} p_k \lambda_l^{(k)}} |v_j^{(i)}\rangle\langle v_k^{(k)}| \otimes |(i,j)\rangle\langle(k,l)| \right) = \\ &= \sum_{m=1}^d \langle m | \otimes \text{id} \left( \sum_{i=1}^N \sum_{j=1}^d \sum_{k=1}^N \sum_{l=1}^d \sqrt{p_i \lambda_j^{(i)} p_k \lambda_l^{(k)}} |v_j^{(i)}\rangle\langle v_k^{(k)}| \otimes |(i,j)\rangle\langle(k,l)| \right) |m \rangle \otimes \text{id} \\ &= \sum_{m=1}^d \sum_{i=1}^N \sum_{j=1}^d \sum_{k=1}^N \sum_{l=1}^d \sqrt{p_i \lambda_j^{(i)} p_k \lambda_l^{(k)}} \langle m | v_j^{(i)} \otimes v_k^{(k)} | m \rangle |(i,j)\rangle\langle(k,l)| \\ &= \sum_{i=1}^N \sum_{j=1}^d \sum_{k=1}^N \sum_{l=1}^d \sqrt{p_i \lambda_j^{(i)} p_k \lambda_l^{(k)}} \langle v_k^{(k)} | v_j^{(i)} |(i,j)\rangle\langle(k,l)| \end{aligned}$$

If I perform a measurement on B (with  $\frac{\text{PV}}{\text{P}_{AB}}$  with  $\langle(i,j)\rangle_{\text{PV}}$ ), then the post measurement state will be:

$$P_B' = \sum_{i=1}^N \sum_{j=1}^d \text{tr}(P_B \rho_{ij}) \left( \frac{P_{ij} P_B P_{ij}}{\text{tr}(P_B \rho_{ij})} \right)$$

$$= \sum_{i=1}^N \sum_{j=1}^d P_{ij} P_B P_{ij} = \sum_{i,j} P_{ij} \lambda_j^{(i)} |(i,j)\rangle\langle(i,j)|$$

$$P_B' = \sum_{i,s} P_i \lambda_s^{(i)} |(i,s)\rangle\langle(i,s)|$$

$$S(P_B') = H\left(\sum_i P_i \lambda_s^{(i)}\right) = H\left(\sum_i P_i\right) + \sum_i P_i \underbrace{\sum_{s=1}^d \lambda_s^{(i)} \log(\lambda_s^{(i)})}_{S(P_i)}$$

\* LEMMA •  $\sum_k P_k$  PVH set  $\Rightarrow S\left(\sum_k P_k p_k\right) \geq S(p)$

(PROOF LATER)

$\uparrow$   
state after measurement.

\*  $S(P_B') \geq S(P_B) = S\left(\sum_i P_i p_i\right)$

$\parallel$   
 $H\left(\sum_i P_i\right) + \sum_i P_i S(P_i)$

\* LEMMA •  $\sum_k P_k$  PVH set  $\Rightarrow S\left(\sum_k P_k p_k\right) \geq S(p)$

$\uparrow$   
state after measurement.

SUB PROOF:

\*  $0 \leq S(p||p') = -S(p) - \text{tr}(p \log(p'))$

KLEIN INEQ

\*  $- \text{tr}(p \log(p')) \geq S(p)$

$$-\text{tr}(P \log(P)) = -\text{tr}(P' \log(P')) = S(P') \Rightarrow \square$$

SUB SUB PROOF:

$$\text{tr}(P \log(P)) = \text{tr}\left(\sum_i P_i p \log(p^i)\right) = \sum_i \text{tr}\left(P_i^2 p \log(p^i)\right) = \sum_i \text{tr}\left(P_i p \log(p^i) P_i\right) =$$

( $P \in \mathbb{R}^{n \times n}$   $\Rightarrow P_i P_j = \delta_{ij} P_i$ )  
 $\sum_i P_i = I$ )

$$\log(p^i) P_i = P_i \log(p^i)$$

$$P' = \sum_j P_j P_j \Rightarrow P' P_i = P_i P_j P_i = P_i P' \Rightarrow \log(p^i) P_i = P_i \log(p^i)$$

$\log(p^i) = \sum_{k=0}^{\infty} c_k i^k$

$$= \sum_i \text{tr}\left(P_i P_j \log(p^i)\right) = \sum_i \text{tr}(P' \log(P')) \Rightarrow \square$$

**THEOR.**  
(STRONG SUB-ADDITIONITY)  
(SSA)

$$S(P_{ABC}) + S(P_C) \leq S(P_{AC}) + S(P_{BC})$$

with  $P_{AB} = \text{tr}_C(P_{ABC})$ ,  $P_{BC} = \text{tr}_A(P_{ABC})$ ,  $P_A = \text{tr}_{BC}(P_{ABC})$ , ...

**OBS** STRONG SUB-ADDITIONITY  $\Rightarrow$  SUB-ADDITIONITY

PROOF:

$$P_{ABC} := P_{AB} \otimes |0\rangle\langle 0|$$

$$\underbrace{S(P_{ABC})}_{\substack{\text{PRODUCT STATE} \\ \rightarrow |1\rangle}} + \underbrace{S(P_C)}_{\substack{|1\rangle \\ S(|0\rangle\langle 0|)}} \leq \underbrace{S(P_{AC})}_{\substack{|1\rangle \\ S(P_A \otimes |0\rangle\langle 0|)}} + \underbrace{S(P_{BC})}_{\substack{|1\rangle \\ S(P_B \otimes |0\rangle\langle 0|)}}$$

$|0\rangle \leftarrow \text{PURE}$

$$\Rightarrow S(P_{AB}) \leq S(P_A) + S(P_B) \Rightarrow \square$$

PROP.

$$\cdot P^{\otimes h} := \underbrace{P \otimes \dots \otimes P}_{n\text{-times}}$$

$$\cdot S(P^{\otimes h}) = \underset{\substack{\uparrow \\ \text{SUB-APP.} \\ (\text{PRODUCT})}}{h S(P)}$$

$$\circ S(P^{\otimes h} \parallel \tau^{\otimes h}) = h S(P \parallel \tau)$$

PROOF:

$$\begin{aligned} S(P^{\otimes h} \parallel \tau^{\otimes h}) &= -S(P^{\otimes h}) + \text{tr}(P^{\otimes h} \log(\tau^{\otimes h})) = \\ &= -h S(P) + h \text{tr}(P \log(\tau)) = h S(P \parallel \tau). \end{aligned}$$

$\left( \begin{array}{l} \cdot S(P^{\otimes h}) = h S(P) \\ \cdot \log(A \otimes B) = \log A \otimes \mathbb{1} + \mathbb{1} \otimes \log(B) \end{array} \right)$